## Extending the swsusp Hibernation Framework to ARM

**Russell Dill** 

#### Introduction

- Russ Dill of Texas Instruments
- swsusp/hibernation on ARM
  - Overview
  - Challenges
  - Implementation
  - Remaining work
  - Debugging
- swsusp restore from U-Boot
- Code at:
  - https://github.com/russdill/linux/commits/arm-hibernation-am33xx
  - https://github.com/russdill/commits/hibernation
  - eLinux.org page: http://elinux.org/ARM\_Hibernation

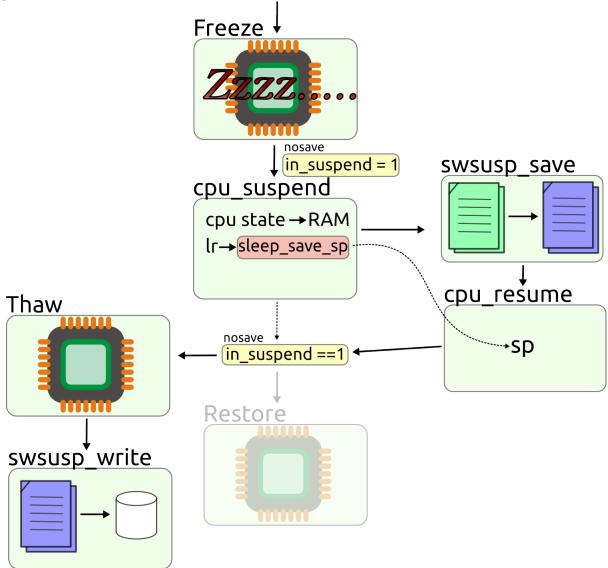
#### **Motivation**

- Hibernation provides zero power consumption sleep
- Allows for snapshot boot
- Shares requirements with self-refresh only sleep modes
  - RTC-Only+DDR self-refresh

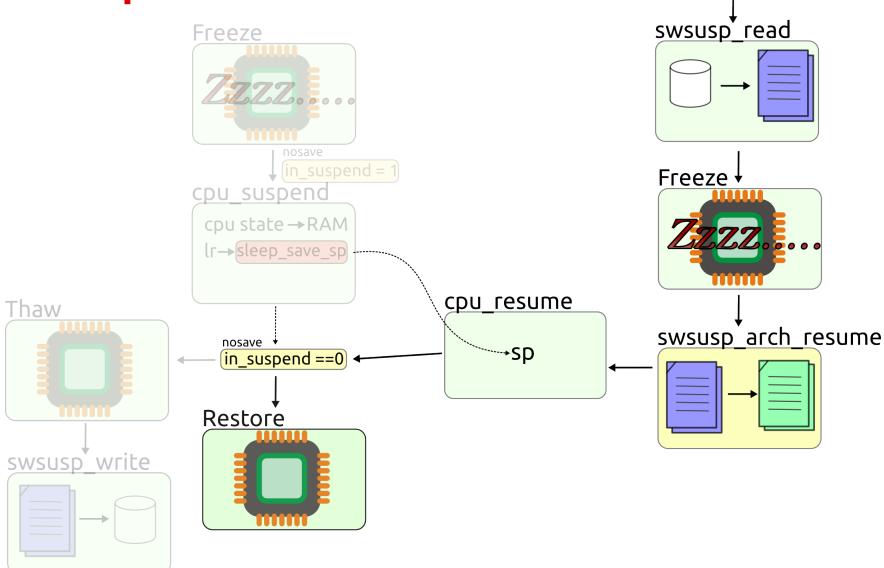
#### swsusp

- Mainline hibernation implementation since 2.6.0
  - TuxOnIce (Suspend2)
- Uses swap device to store image
- Requires 1/2 of system RAM to be free
- Can be used with uswsusp to support additional features
  - Encryption
  - Limitless storage options
  - Graphical progress

#### swsusp

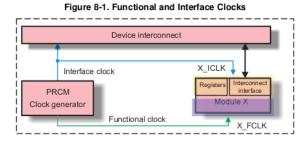


#### **swsusp**



#### **OMAP PM**

- Clocks
  - Clock gating
  - Clock domains
  - Clock scaling
- Power
  - Power domains
    - Logic
    - Retention
  - Voltage scaling
- PRCM Controls these features



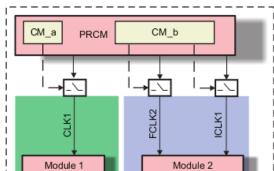
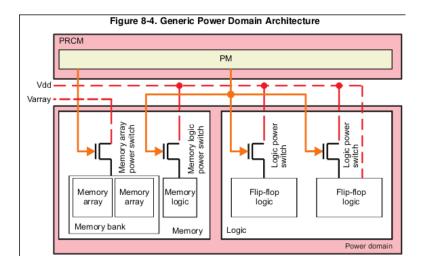


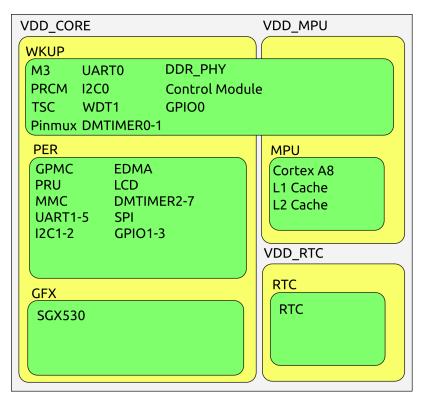
Figure 8-2. Generic Clock Domain





#### **AM33xx PM Overview**

- MPU, PER, and GFX power domains can be turned off during suspend
- Current OMAP PM core assumes
   WKUP domain will always have
   power



#### **WKUP Context**

- Used for:
  - Power, reset, and clock management (PRCM)
  - Pin mux configuration
  - modules that wake up the processor from suspend
    - Such as UART0, GPIO0
  - modules that should continue running when processor is in suspend
    - Such as DMTIMER0
- After hibernation, we need to restore this state
  - Many critical portions of kernel code depend on devices within the WKUP domain
    - sched\_timeout

#### **PRCM – Power Domains**

- Represented by arch/arm/mach-omap2/powerdomain.c
  - Power state configuration

#### PRCM - Reset State/Module State

Represented by omap\_hwmod, leverage it

#### **PRCM – Clock Domains**

Represented by arch/arm/mach-omap2/clockdomain.c

```
static int am33xx_clkdm_save_context(struct clockdomain *clkdm)
        clkdm->context = am33xx_cm_read_reg_bits(clkdm->cm_inst,
                                clkdm->clkdm_offs, AM33XX_CLKTRCTRL MASK);
        return 0;
static int am33xx clkdm restore context(struct clockdomain *clkdm)
        switch (clkdm->context) {
        case OMAP34XX_CLKSTCTRL_DISABLE_AUTO:
                am33xx clkdm deny idle(clkdm);
                break;
        case OMAP34XX CLKSTCTRL FORCE SLEEP:
                am33xx_clkdm_sleep(clkdm);
                break:
        case OMAP34XX CLKSTCTRL FORCE WAKEUP:
                am33xx_clkdm_wakeup(clkdm);
                break;
        case OMAP34XX CLKSTCTRL ENABLE AUTO:
                am33xx clkdm allow idle(clkdm);
                break;
        return 0;
```

#### **PRCM - Clocks**

Leverage the clock tree by adding context save/restore callbacks

```
static int clk_save_context(struct clk *clk)
        struct clk *child;
        struct hlist_node *tmp;
        int ret = 0;
        hlist_for_each_entry(child, tmp, &clk->children, child_node) {
                ret = clk_save_context(child);
                if (ret < 0)
                        return ret;
        }
        if (clk->ops && clk->ops->save_context)
                ret = clk->ops->save_context(clk->hw);
        return ret;
static void clk restore context(struct clk *clk)
{
        struct clk *child;
        struct hlist_node *tmp;
        if (clk->ops && clk->ops->restore_context)
                clk->ops->restore_context(clk->hw);
        hlist_for_each_entry(child, tmp, &clk->children, child_node)
                clk_restore_context(child);
```

- Controls how internal signals are routed to external pins
- Contains memory map of register area, but no complete description of registers
- AM335X errata complicates the situation, certain registers lose context when the PER domain powers during suspend
- The pinctrl subsystem needs knowledge of which registers are available, and which domain they are in.

Temporary measure, list each power domain register set as a pinconf function

```
am33xx pinmux: pinmux@44e10800 {
        compatible = "pinctrl-single";
        reg = <0x44e10800 0x0238>;
        #address-cells = <1>;
       #size-cells = <0>;
       pinctrl-single,register-width = <32>;
        pinctrl-single,function-mask = <0x7f>;
       pinctrl-names = "default", "context";
        pinctrl-0 = <>;
        pinctrl-1 = <&am33xx pmx wkup &am33xx pmx per>;
       am33xx pmx wkup: am33xx pmx wkup {
                pinctrl-single,pins = <</pre>
                        0x000 0x00 /* GPMC AD0 */
                        0x004 0x00 /* GPMC_AD1 */
                        0x008 0x00 /* GPMC_AD2 */
                        0x00c 0x00 /* GPMC AD3 */
                        0x010 0x00 /* GPMC AD4 */
                        0x014 0x00 /* GPMC AD5 */
                        0x1f8 0x00 /* RTC_PWRONRSTN */
                       0x1fc 0x00 /* PMIC POWER EN */
                       0x200 0x00 /* EXT WAKEUP */
                       0x204 0x00 /* RTC KALDO ENN */
                       0x21c 0x00 /* USB0 DRVVBUS */
                        0x234 0x00 /* USB1 DRVVBUS */
               >;
       };
```

```
am33xx_pmx_per: am33xx_pmx_per {
                 pinctrl-single,pins = <</pre>
                        0x040 0x00 /* GPMC_A0 */
                         0x044 0x00 /* GPMC_A1 */
                         0x048 0x00 /* GPMC A2 */
                        0x04c 0x00 /* GPMC A3 */
                         0x050 0x00 /* GPMC A4 */
                         0x054 0x00 /* GPMC A5 */
                         0x058 0x00 /* GPMC A6 */
                        0x134 0x00 /* MII1 RXD3 */
                         0x138 0x00 /* MII1_RXD2 */
                         0x13c 0x00 /* MII1 RXD1 */
                        0x140 0x00 /* MII1_RXD0 */
                        0x144 0x00 /* MII1_REFCLK */
                        0x148 0x00 /* MDIO DATA */
                        0x14c 0x00 /* MDIO CLK */
                >;
        };
};
```

Code added to pinctrl to save/restore a pinctrl function group

```
int pinmux_save_context(struct pinctrl_dev *pctldev, const char *function)
        const struct pinmux ops *pmxops = pctldev->desc->pmxops;
        ret = pinmux func name to selector(pctldev, function);
        if (ret < 0) {</pre>
                dev_err(pctldev->dev, "invalid function %s\n", function);
                return ret;
        }
        if (!pmxops || !pmxops->save_context)
                return -EINVAL;
        return pmxops->save_context(pctldev, ret);
EXPORT SYMBOL(pinmux save context);
void pinmux_restore_context(struct pinctrl_dev *pctldev, const char *function)
        const struct pinmux_ops *pmxops = pctldev->desc->pmxops;
        int ret;
        ret = pinmux_func_name_to_selector(pctldev, function);
        if (ret < 0) {
                dev_err(pctldev->dev, "invalid function %s\n", function);
                return:
        }
        if (!pmxops || !pmxops->restore context)
                return;
        pmxops->restore_context(pctldev, ret);
EXPORT_SYMBOL(pinmux_restore_context);
```

```
static int pcs_save_context(struct pinctrl_dev *pctldev, unsigned fselector)
        struct pcs_device *pcs;
        struct pcs_function *func;
        pcs = pinctrl_dev_get_drvdata(pctldev);
        func = radix tree lookup(&pcs->ftree, fselector);
        if (!func) {
                dev_err(pcs->dev, "%s could not find function%i\n",
                         _func__, fselector);
                return -ENODEV;
        }
        for (i = 0; i < func->nvals; i++) {
                struct pcs_func_vals *vals;
                vals = &func->vals[i];
                vals->val = pcs->read(vals->reg);
        }
        return 0:
}
static void pcs_restore_context(struct pinctrl_dev *pctldev, unsigned fselector)
        struct pcs_device *pcs;
        struct pcs_function *func;
        int i;
        pcs = pinctrl_dev_get_drvdata(pctldev);
        func = radix_tree_lookup(&pcs->ftree, fselector);
        if (!func) {
                dev_err(pcs->dev, "%s could not find function%i\n",
                        __func__, fselector);
                return:
        for (i = 0; i < func->nvals; i++) {
                struct pcs_func_vals *vals;
                unsigned val, mask;
                vals = &func->vals[i];
                val = pcs->read(vals->reg);
                if (!vals->mask)
                        mask = pcs->fmask;
                else
                        mask = pcs->fmask & vals->mask;
                val &= ~mask;
                val |= (vals->val & mask);
                pcs->write(val, vals->reg);
```



- Current solution is a bit of a hack and likely not upstreamable.
- Possible solution?
  - New type of pinctrl register grouping
  - Would contain reference to power domain register group is contained in
  - Code could use syscore suspend/resume callbacks to save and restore context
- Problem
  - omap2+ power domains are currently arch specific

#### clocksource/clockevent

- clockevent is already handled properly, disabling on suspend and reprogramming on resume
  - Located in PER power domain
- clocksource is assumed to be always running and within a domain that does not lose power
- clocksource is also required for many kernel delay calculations. Must be restored before most other kernel code

```
static cycle_t clksrc_suspend_cyc;
                                                                                 static void omap_clksrc_resume(struct clocksource *cs)
static void omap_clksrc_suspend(struct clocksource *cs)
                                                                                         char name[10];
                                                                                         struct omap_hwmod *oh;
        char name[10];
                                                                                         u32 ctx_loss_cnt_after;
        struct omap_hwmod *oh;
                                                                                         sprintf(name, "timer%d", clksrc.id);
       sprintf(name, "timer%d", clksrc.id);
                                                                                         oh = omap_hwmod_lookup(name);
        oh = omap_hwmod_lookup(name);
                                                                                         if (!oh)
       if (!oh)
                                                                                                  return:
               return;
                                                                                         ctx_loss_cnt_after = omap_hwmod_get_context_loss_count(oh);
        clksrc_suspend_cyc = (cycle_t) _ omap_dm_timer_read_counter(&clksrc, 1);
                                                                                         if (ctx_loss_cnt_after != clksrc.ctx_loss_count) {
        clksrc.ctx_loss_count = omap_hwmod_get_context_loss_count(oh);
                                                                                                  omap2_dflt_clk_restore_context(__clk_get_hw(clksrc.fclk));
                                                                                                  omap_hwmod_reset(oh);
                                                                                                  __omap_dm_timer_load_start(&clksrc,
                                                                                                                  OMAP_TIMER_CTRL_ST | OMAP_TIMER_CTRL_AR,
                                                                                                                  clksrc_suspend_cyc, 1);
                                                                                                  __omap_dm_timer_int_enable(&clksrc, OMAP_TIMER_INT_OVERFLOW);
                                                                                 }
```

#### **SRAM**

- Internal memory on many OMAP processors used to run suspend resume code or code that modifies memory controller registers or clocking
- Currently restored only for OMAP3, but in an OMAP3 specific way –
   Make it more general instead

- Many drivers just need to know that their power domain lost context
- Teach arch/arm/mach-omap2/powerdomain.c about hibernation induced off modes.

- Many drivers that depend on a context loss count function pointer do not get that pointer under DT based systems
  - gpio-omap
  - omap\_hsmmc
  - omap-serial
- Currently a hack fix with a pointer to omap\_pm\_get\_dev\_context\_loss\_count
  - [HACK]: Crosses arch/arm/mach-omap2, drivers/ boundary
- There is a need for a generic framework to inform drivers when their device has lost power

- Some drivers are misconfigured in such a way to prevent
  - suspend/resume callbacks during hibernation
- When not using dev\_pm\_ops, the platform\_driver .suspend/.resume callbacks are used for hibernation thaw/freeze/restore/poweroff functionality
- However, when using dev\_pm\_ops these must be filled in. The helper

```
.owner = THIS MODULE,
                                             .pm = &omap hsmmc dev pm ops,
                                             .of match table = of match ptr(omap mmc of match),
                                          },
macro, SET_SYSTEM_SLEEP_PM_OPS should be used to fill in the
thaw/freeze/restore/poweroff callbacks (unless special
thaw/freeze/restore/poweroff behavior is required).
```

static struct dev pm ops omap hsmmc dev pm ops = { = omap hsmmc suspend,

= omap\_hsmmc\_prepare,

= omap\_hsmmc\_complete,

.runtime suspend = omap hsmmc runtime suspend, .runtime resume = omap hsmmc runtime resume,

static struct platform\_driver omap\_hsmmc\_driver = {

= omap hsmmc probe, = omap hsmmc remove,

= {

.name = DRIVER NAME,

= omap hsmmc resume,

.suspend

.resume

.prepare

.complete

.probe

.remove

.driver

- Some drivers \*do\* need special hibernation callbacks
- The omap watchdog requires special handling because the state of the watchdog under the boot kernel is not known

```
static int omap_wdt_restore(struct device *dev)
        struct watchdog device *wdog = dev get drvdata(dev);
       struct omap wdt dev *wdev = watchdog get drvdata(wdog);
       omap wdt resume(dev);
         * We don't know what the resume kernel last pinged the WDT with. If
         * it pinged it with the same value we ping it with, the ping will be
         * ignored. Double ping to be sure we reset the timer.
       if (wdev->omap_wdt_users)
               omap wdt ping(wdog);
       return 0;
static const struct dev pm ops omap wdt pm ops = {
        .suspend
                       = omap wdt suspend,
        .freeze
                       = omap_wdt_suspend,
        .poweroff
                       = omap wdt suspend,
                       = omap wdt resume,
        .resume
                       = omap wdt resume,
        .thaw
                       = omap wdt restore,
        .restore
};
```

## Saving/Restoring WKUP Domain

Putting it all together in pm33xx.c

```
static int am33xx_wkup_save_context(void)
        int ret;
        ret = pinmux_save_context(pmx_dev, "am33xx_pmx_wkup");
        if (ret < 0)
                return ret;
        omap into save context();
        am33xx control save context();
        clks save context();
        pwrdms_save_context();
        omap hwmods save context();
        clkdm save context();
        omap_sram_save_context();
        return 0;
static void am33xx wkup restore context(void)
        clks_restore_context();
        pwrdms_restore_context();
        clkdm_restore_context();
        omap hwmods restore context();
        am33xx control restore context();
        pinmux restore context(pmx dev, "am33xx pmx wkup");
        omap intc restore context();
        wkup_m3_reinitialize();
        omap_sram_restore_context();
```

## **Generic ARM Hibernation Support**

## **Hibernation support for ARM**

- Minimum implementation
  - swsusp\_arch\_suspend
    - Save current cpu state
    - Call swsusp\_save to snapshot memory
    - Return control to swsusp\_arch\_suspend caller
  - swsusp\_arch\_resume
    - Perform page copies of pages in the restore\_pbelist
      - Copies pages from their restored location to their original location
    - Restore cpu state from swsusp\_arch\_suspend
    - · Return control to swsusp arch suspend caller
  - pfn\_is\_no\_save
    - Return true if this pfn is not to be saved in the hibernation image
  - save\_processor\_state (Called just before swsusp\_arch\_suspend)
    - Save any extra processor state (fp registers, etc)
  - restore\_processor\_state (Called just after swsusp\_arch\_suspend)
    - Restore extra processor state
- Based on a patchset from Frank Hofmann



## Hibernation support - arch\_suspend

- swsusp\_arch\_suspend
  - Utilizes cpu suspend to save current cpu state
  - Second argument of cpu\_suspend is called after state is saved
  - Calling cpu\_resume
     causes execution to
     return to cpu\_suspend
     caller
  - Utilizing soft\_restart disables MMU as cpu\_resume expects

```
* Snapshot kernel memory and reset the system.
 * After resume, the hibernation snapshot is written out.
static int notrace swsusp arch save image(unsigned long unused)
        extern int swsusp save(void);
       int ret;
        ret = swsusp_save();
       if (ret == 0)
                soft restart(virt to phys(cpu resume));
        return ret:
 * Save the current CPU state before suspend / poweroff.
int notrace swsusp_arch_suspend(void)
        return cpu_suspend(0, __swsusp_arch_save_image);
```

## Hibernation support - arch\_resume

- swsusp\_arch\_resume
  - Uses stack allocated in nosave region to prevent ourselves from overwriting our stack
  - We will overwrite our code, but with the same bytes
  - Uses cpu\_resume to restore cpu state and return to cpu\_suspend caller (swsusp\_arch\_suspend)

```
* The framework loads the hibernation image into a linked list anchored
 * at restore_pblist, for swsusp_arch_resume() to copy back to the proper
 * destinations.
 * To make this work if resume is triggered from initramfs, the
 * pagetables need to be switched to allow writes to kernel mem.
static void notrace swsusp arch restore image(void *unused)
        extern struct pbe *restore pblist;
        struct pbe *pbe;
        cpu_switch_mm(idmap_pgd, &init_mm);
        for (pbe = restore pblist; pbe; pbe = pbe->next)
                copy page(pbe->orig address, pbe->address);
        soft_restart_noirq(virt_to_phys(cpu_resume));
static u8 __swsusp_resume_stk[PAGE_SIZE/2] __nosavedata;
 * Resume from the hibernation image.
 * Due to the kernel heap / data restore, stack contents change underneath
 * and that would make function calls impossible; switch to a temporary
 * stack within the nosave region to avoid that problem.
int __naked swsusp_arch_resume(void)
        extern void call with stack(void (*fn)(void *), void *arg, void *sp);
        cpu_init(); /* get a clean PSR */
        call_with_stack(__swsusp_arch_restore_image, 0,
                __swsusp_resume_stk + sizeof(__swsusp_resume_stk));
        return 0;
```

## **AM33xx Hibernation Support**

- With prep work done, adding hibernation support to AM33xx is actually fairly straightforward
- begin/end wrap all hibernation code
- We use disable/enable\_hlt to prevent pm\_idle from being called

- The enter call back just powers down the machine
- These calls make sure that the hardware is in the same state before running the restored image as when it was made

```
static int am33xx_hibernation_begin(void)
{
         disable_hlt();
         return 0;
}
static void am33xx_hibernation_end(void)
{
         enable_hlt();
}
```

```
static int am33xx_hibernation_enter(void)
{
    machine_power_off();
    return 0;
}

static int am33xx_hibernation_pre_restore(void)
{
    omap2_gpio_prepare_for_idle(1);
    return 0;
}

static void am33xx_hibernation_restore_cleanup(void)
{
    omap2_gpio_resume_after_idle();
}
```



## **AM33xx Hibernation Support**

- pre\_snapshot saves all our state registers and prepares the GPIOs for power loss
- leave is called after restoring an image. We inform the power domains that they have lost power and we restore our wkup context
- restoring an image (after leave) and after snapshotting the system. We continue our context restore and also undo the actions in pre\_snapshot

```
static int am33xx_hibernation_pre_snapshot(void)
        am33xx_per_save_context();
        omap2 gpio prepare for idle(1);
        am33xx wkup save context();
        return 0;
static void am33xx hibernation leave(void)
        pwrdms lost power();
        am33xx wkup restore context();
static void am33xx hibernation_finish(void)
        omap2 gpio resume after idle();
        am33xx_per_restore_context();
```

## **Debugging Methods**

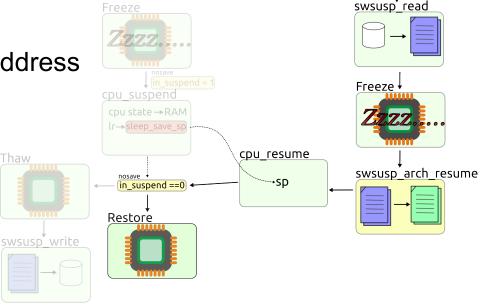
- Debugging can be difficult as the hardware is usually in some unknown state.
- Debugging using GPIOs
  - GPIOs are usually pretty easy to configure clocks for and enable with just a few register writes, even from assembly
  - Binary search of where the code is failing can be performed by moving the GPIO enable around
- printk
  - The kernel logging facility is useful so long as you are getting to a point where serial output is enabled
- openocd
  - Usually PC is down in the weeds
  - openocd+gdb?
- Register map comparisons
  - Utilizing devmem2 to snapshot register values before and after a hibernation file is useful to track down missed registers or buggy restore code

#### **Restore from U-Boot**



#### swsusp and U-Boot

- Restoring from hibernation just involves copying pages from disk into memory and jumping to an address
  - Thats what U-Boot does!
- Restoring from U-Boot can be faster than booting a kernel just to copy pages
- Issues
  - U-Boot has no idea what address to jump to
  - U-Boot doesn't know the contents or even location of the nosave pages



#### **Kernel Modifications – nosave section**

- U-Boot doesn't know about nosave pages or their address
- We instead save and restore them from the kernel
- Backup nosave pages are saved at boot
- Special version of cpu\_resume is provided that restores nosave pages before calling the real cpu\_resume

```
static int __init swsusp_arch_init(void)
{
    char *backup;
    size_t len;

    len = &__nosave_end - &__nosave_begin;
    backup = kmalloc(len, GFP_KERNEL);
    if (backup) {
            pr_info("%s: Backed up %d byte nosave region\n", __func__, len);
            memcpy(backup, &__nosave_begin, len);
    }

    __nosave_backup_phys = virt_to_phys(backup);
    __nosave_begin_phys = virt_to_phys(&__nosave_begin);
    __nosave_end_phys = virt_to_phys(&__nosave_end);

    return 0;
}
late_initcall(swsusp_arch_init);
```

```
ENTRY(cpu resume restore nosave)
        ldr
                r0, = nosave backup phys
        ldr
                r0, [r0]
                r1, =__nosave_begin_phys
        ldr
        ldr
                r1, [r1]
                r2, =__nosave_end_phys
        ldr
        ldr
                r2, [r2]
        ldmia
                r0!, {r3-r10}
1:
                r1!, {r3-r10}
        stmia
                г1, г2
        CMD
                1b
        bne
                cpu_resume
```

## Kernel Modifications - cpu\_resume

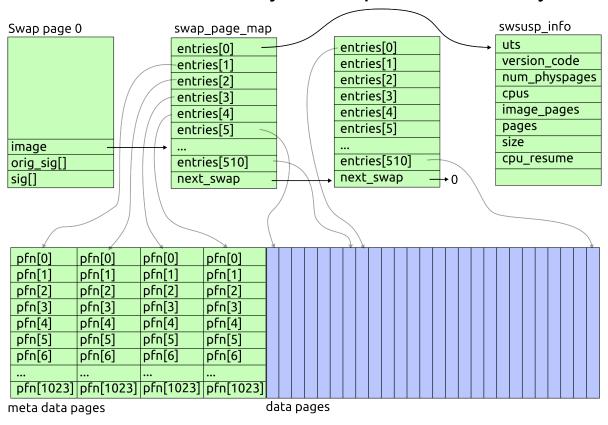
- Need to pass address of cpu\_resume function to U-Boot
  - Store in swsusp\_info page
  - Add arch callback for storing that data in the swsusp\_info page
- Just stores the physical address of the new version of cpu\_resume that first copies the nosave pages

## swsusp Image Layout

Each metadata entry is associated with the same numbered data page

Each data page is to be loaded into memory at the pfn indicated by

its metadata pfn entry



#### **U-Boot modifications**

- Provide cmd\_swsusp
  - No-op if S1SUSPEND sig does not exist

```
U-Boot# help swsusp
swsusp - Restore SWSUSP hibernation image

Usage:
swsusp <interface> [<dev[:part]>] [<offset>]
```

- Rewrites sig with orig\_sig to prevent boot loop on bad image
  - Snapshot booting can populate orig\_sig with S1SUSPEND
- Reads in metadata pages with pfn mappings
  - Also populates bitmap of used pages for easy access to free pages
- Copy each data page to memory
  - Original location if it is free
  - Other wise copy to first available free page and update remap list
- Copy finish function and cpu\_resume address to free data page
- Run finish function from free data page (use stack contained in free page)
  - Copies remapped pages to their correct location
  - Jumps to cpu\_resume function



## **U-Boot Memory Mapping**

- The U-Boot memory mapping makes it very easy to see if we can load a page directly into its original location
- If not, we load it into a location not used by U-Boot or the final location of any of the swsusp pages

Vectors

Free

Stack
Global data
Board info
Malloc space
U-Boot data & bss

LCD

## Loading pfn and Free Page Mapping

- We utilize malloc'd pages to store the pfn index
- Mark used pages as we go

## **Loading swsusp Pages Into Memory**

- Utilize free pages to store remapping lists, malloc'd data will be overwritten
- min\_page is first free page in U-Boot memory map
- max\_page is last free page in U-Boot memory map (well before stack pointer)
- If a page is to be copied into U-Boot's memory space, it is instead copied into an unused free page

## **Prepare to Copy Remapped Pages**

- Final copy must happen from memory unused by swsusp or U-Boot
  - remap\_orig/remap\_temp already exist in free page
  - Utilize free page for final copy of remapped pages
    - Copy swsusp\_finish into page
    - Copy context information into page
    - Setup stack pointer at end of page

```
/* put end markers on the remap list */
remap orig[remap idx] = (void *) ~OUL;
remap_temp[remap_idx] = (void *) ~OUL;
/* Make a copy of swsusp_finish in a free data page */
data page = pg2addr(free page get next());
memcpy(data page, swsusp finish, PAGE SIZE);
swsusp finish copy = (void *) data page;
/* Setup context for swsusp finish */
context = (struct swsusp finish context *) (data page + PAGE SIZE);
context--:
context->remap_orig = remap_orig;
context->remap temp = remap temp;
context->cpu resume = swsusp info->cpu resume;
/* Get a stack pointer for swsusp finish */
stack addr = ((char *) context) + PAGE SIZE - sizeof(u32);
cleanup before linux();
/* Copy the final data from a safe place */
call with stack(swsusp_finish_copy, context, stack_addr);
```

## **Copy Remaining Pages**

- Moved remapped pages into their originally intended location
- Call cpu\_resume (actually cpu\_resume\_copy\_nosave)

```
static void swsusp finish (void *userdata)
        struct swsusp_finish_context *context = userdata;
        while (*context->remap_orig != (void *) ~OUL) {
                u32 *orig, *temp;
                int count;
                count = PAGE SIZE / 4;
                orig = *context->remap orig;
                temp = *context->remap_temp;
                while (count--)
                        *orig++ = *temp++;
                context->remap orig++;
                context->remap temp++;
        context->cpu resume();
}
```

## **Lessons Learned/Open Issues**

- Save/Restore context function
  - Bit of a hack, better to keep context as we go and look at context loss count
    - /dev/mem bitbangers?
    - Registers that are initialized at boot and never read by the driver?
- Over zealous restore
  - Don't restore context of clocks who's parents are disabled
    - Check to make sure clock drivers are programming appropriate registers on resume
- Setup debug framework early
- pinctrl
  - Properly integrate with system

## **Questions?**

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